RESEARCH ARTICLE   OPEN ACCESS

A Review of Energy Efficient In-network Data Aggregation Protocols

Amit S. Bhosale *, Jaishri M. Waghmare **, Manisha S. Mahindrakar***

*(Department of Computer Science and Engineering, SRTMU University, Nanded-431606
**(Department of Computer Science and Engineering, SRTMU University, Nanded-431606
**(Department of Computer Science and Engineering, SRTMU University, Nanded-431606

ABSTRACT
A Wireless Sensor Network (WSN) consists of hundreds or thousands of sensor nodes. The basic role of these nodes is to collect data from an environment and send it to sink node. This process generates huge amount of data. The communication among sensor nodes causes high power consumption. The main purpose of in-network data aggregation protocols is to collect and aggregate data in an energy efficient manner so that network lifetime is improved. This paper provides a review of existing literature on techniques and protocols for in-network data aggregation in WSN.
Keywords - Data aggregation, Energy Saving, Wireless Sensor Network.

I. Introduction

1.1. Wireless Sensor Network

A WSN consists of scattered autonomous sensors to observe physical or environmental conditions (such as pressure, temperature, and sound etc.) and to send their data through the network to a base station. WSNs are deployed in security applications such as factory monitoring, environmental monitoring, burglar alarms and fire alarms. The sensor nodes for these applications are typically deployed in unsecured locations and are not made tamper-proof due to cost considerations. Hence, an adversary or an attacker could take control of one or more sensor nodes and launch active attacks to threaten correct network operations. Such environments pose a particularly challenging set of constraints for the protocol designer. Sensor network protocols must be highly energy efficient while being able to function securely in the presence of possible malicious nodes within the network [1].

Sensor networks made up of small self-computing devices capable of producing digital representations of real-world environment phenomena. Due to size and battery power limitations, these devices typically have limited storage capacity, limited energy resources, and limited network bandwidth. Data produced by nodes in the network propagates through the network via wireless links. When compared to local processing of data, wireless transmission is extremely expensive. Researchers at the University of California, Santa Barbara estimate that sending a single bit over radio is at least three orders of magnitude more expensive than executing a single instruction. The limited amount of energy, bandwidth, and storage capacity available to sensor nodes calls for specialized optimizations of queries injected into the network.

Initially, in-network aggregation techniques involved different ways to route packets in order to combine data coming from different sources but directed towards the same destination. These protocols were simply routing algorithms which differed from more traditional ad hoc routing protocols in the metric they used to select the routing paths [3]. In-network data aggregation is a complex problem that involves many layers of the protocol stack and different aspects of protocol design, and a characterization and classification of concepts [2].

The remainder of this paper is organized as follows. Section 2 introduces about the In-Network Aggregation. Protocols in In-Network data aggregation are categorized in tree based, clusters based and structure less approaches and their details in given in section 3. Finally, section 4 concludes the paper with a comparative summary of the surveyed approaches.

1.2. In-network Aggregation

"In-network aggregation is the global process of gathering and routing information through a multihop network, processing data at intermediate nodes with the objective of reducing resource consumption (in particular energy), thereby increasing network lifetime" [3].

We can distinguish between two approaches:
- With size reduction

It refers to the process of merging and reducing data coming from different sources in order to reduce the information to be sent over the network. As an example, assume that a node receives two

packets from two different sources containing the locally measured temperatures. Instead of forwarding the two packets, the sensor may compute the average of the two readings and send it in a single packet.

- **Without size reduction**
  It refers to the process of merging packets coming from different sources into the same packet without data processing: assume to receive two packets carrying different physical quantities, e.g. temperature and humidity. These two values cannot be processed together but they can still be transmitted in a single packet, thereby reducing overhead.

II.   Existing In-network Data Aggregation Protocols

Data aggregation process is performed by specific routing protocol. The main aim is aggregating data to minimize the energy consumption. So sensor nodes should route packets based on the data packet content and choose the next hoping order to promote in-network aggregation. Basically routing protocol is divided by the network structure, that's why routing protocols in in-network data aggregation are categorized in tree based, clusters based and structure less approaches.

2.1 Tree Based Protocols

The main aim of tree based protocols is to maintain the energy consumption of the sensor nodes by using multihop communication within the cluster and by performing data aggregation in order to reduce the number of transmitted messages. Tree based approaches includes Direct Diffusion (DD), Shortest Path Tree (SPT) and Center at Nearest Protocol (CNS).

2.1.1   Direct Diffusion

Intanagonwiwat el. al. [5] have developed Direct Diffusion (DD), which consists of several key elements such as interests, data messages, gradients, and reinforcements as shown in figure 2.1. An interest message is a query which states what a user wants. Each interest contains a narration of a sensing task that is supported by a sensor network for collecting data. Typically, sensor network data are collected information of a physical phenomenon. Such data can be an event, which is a short description of the sensed phenomenon. In directed diffusion, data are named using attribute-value pairs. A sensing task is dispersed throughout the sensor network as an interest for naming data. These scattering sets up gradients within the network designed to "draw" events i.e., data matching the interest. Specifically, a gradient is direction state created on each node that receives an interest. The

gradient direction is set to the neighboring node from which the interest is received. Events start flowing toward the originators of interests along multiple gradient paths.
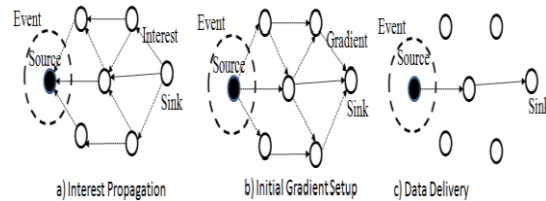

Fig. 2.1 Direct Diffusion

2.1.2   SPT & CNS

B. Krishnamchari et al. in [6] discussed both these protocols assumes the single network flow model which means that there is only one sink node which collects information from the number of data sources. They use a data centric routing scheme which means that during the process of transferring data from sources to the sink the routing nodes can look at the content of data and perform aggregation. In SPT protocol data aggregation scheme all nodes send their information to the sink along the shortest path between the two, and overlapping paths are combined to form the aggregation tree. In CNS protocol all source node sends their data directly to the source which nearest to the sink which sends aggregated data to the sink. The main factor that can affect the process of data aggregation is the placement of source nodes in the network. For this investigation author proposes two models of source placement i.e. the event radius (ER) model and the random source (RS) model. In both models sensor network is generated by scattering n sensor nodes, one of which is sink, in a unit square. In ER model all sources are spread within a range of distance S from random event location and in RS model k nodes are randomly selected to be source. All nodes can communicate with the other nodes in the range of the communication radius R.

2.2 Cluster Based Protocol

Cluster based protocol are similar to the hierarchical organization of the network which includes different protocols like Low Energy Adaptive Clustering Hierarchy (LEACH), Information Based Role Assignment (InFRA), Data Aggregation Aware Routing Protocol (DAARP), and Dynamic Data-Aggregation Aware Routing Protocol (DDAARP), dYnamic and scalablE tree Aware of Spatial Correlation (YEAST).

2.2.1   Leach

Heinzelman, et al. [4] were first to introduce LEACH. It is a self-organizing, adaptive clustering protocol, which includes dynamic and distributed

cluster formation. The diagrammatic representation of LEACH is as shown in figure 2.2. LEACH randomly selects a sensor node as cluster heads (CHs) and switches this role to evenly allocate the energy load between the sensors in the network. In LEACH, the cluster head (CH) nodes merge data received from nodes inside the cluster, and send an aggregated packet to the sink node in order to decrease the amount of information that must be sent to the base station. In LEACH Time Division Multiple Access (TDMA) schedule is used to send data from node to head cluster. Head cluster aggregated data received from node inside the cluster. Communication is via Direct Sequence Spread Sequence (DSSS) and each cluster uses a unique spreading code to reduce inter cluster interference. Also the data from cluster head is sent to the Base Station with the help of unique spreading code and Carrier Sense Multiple Access (CSMA). In LEACH, a good property is to save energy by using aperiodic communication. After suitable time, a randomized rotation of the role of the CH is conducted so that uniform energy dissipation in the sensor network is obtained. The authors found, based on their simulation model that only 5% of the nodes need to act as cluster heads.
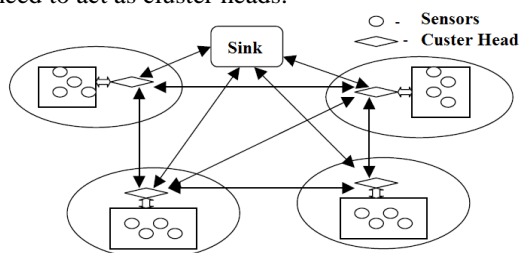

Fig. 2.2　　LEACH

2.2.2　InFRA

Nakamura et al. in [7] discuss the reactive algorithm Information Fusion based Role Assignment (InFRA), in which roles are assigned when any event takes place. Different roles are assign in this protocol like Sink, Collaborator, Coordinator and Relay. In this protocol when multiple nodes detect the same event, they organize themselves into clusters. Then cluster head aggregates data from all cluster members and sends event data towards the sink in multihop fashion. InFRA finds minimum shortest path tree connecting all source nodes to sink such that the intra cluster data aggregation is possible. InFRA provides role migration policy i.e. role of coordinator is transferred from one node to another so that the load of energy consumption is distributed evenly in the nodes inside the cluster. The two types of data aggregation schemes are followed in InFRA i.e. intra-cluster and inter-cluster. In first the data from collaborators are aggregated and in later the data

from coordinator nodes are aggregated. A disadvantage about InFRA algorithm is that for every new event that detect by source nodes the information about that event is broadcasted throughout the network to inform other nodes in the network about its occurrence and to update the paths from the already existing cluster heads to the sink node and events are static in nature and of fixed radius. This limits InFRA scalability.

2.2.3　Daarp

L. A. Villas et al. in [8] overcome the disadvantage of InFRA algorithm. Data Aggregation Aware Routing Protocol (DAARP) is new reliable Data Aggregation Aware Routing Protocol for WSN. Similar to InFRA, in this for each event this algorithm performs the clustering of nodes that detected the event in the network and also the election of cluster head. Then cluster head merges inter cluster data and sends the result to the sink node. After the cluster head formation routes are formed by selecting nodes in the network to existing routing path in which the node in existing path act as an aggregation point. This protocol reduces the number of messages during the setup phase of a routing tree and maximizes overlapping routes. It selects routes with the highest aggregation rate and performs reliable data aggregation transmission and uses fewer control packets to build the paths. Different from InFRA, DAARP does not broadcast a message to the whole network whenever a new event occurs. DAARP is not feasible for scenarios with long duration events because the routes are static, which quickly consumes the energy of the nodes that are part of the routing structure.

2.2.4　Ddaarp

The up gradation of DAARP protocol is discussed by L. A. Villas in [9]. DDAARP builds dynamic routes. In this protocol the routes are configured and processed at the sink node and routes which are created are not dependent on the order of events. Routes established are not kept fixed in DDAARP throughout the duration of events i.e. routes are dynamic in nature. This protocol has low cost in terms of packet control, improves the quality of the routing tree and maximizes information fusion along the routing path. The drawback of this proposal is that packets containing information from nodes tend to increase their size at the information collection stage and this solution becomes impractical for large scale networks. In addition, the sink node needs to have a global knowledge of the network such as node positions, residual energy of nodes and nodes that detected events.

2.2.5 Yeast

L. A. Villas et al. proposes the spatial correlation mechanism in [10]. YEAST is spatial correlation algorithm takes the advantage of best WSN techniques to perform efficient data collection in WSN. It uses spatial correlation model in which a correlation region is defined by the author. In this correlation region the value sensed by the sensor nodes are assumed to be similar and only one sensor node value can represent that value of the whole correlation region. The size of correlation region can be changed dynamically in YEAST. As shown in figure 2.3 the event is divided into cells. Each cell defines a correlated region and only one node within the cell notifies the sensed information. That node is called the leader node or coordinator node. To divide the energy flow evenly within the cell the role of leader transferred from one node to another.
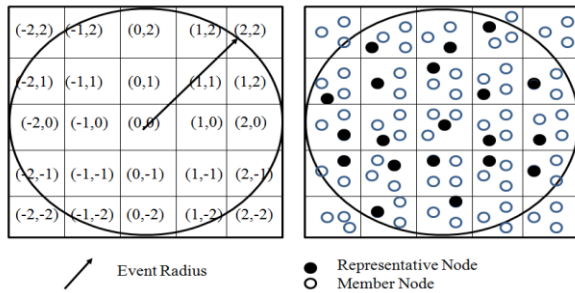

Fig. 2.3 YEAST

2.3 Structureless Protocols

This category includes protocols named as Data Aware Anycast (DAA) [11] and Dynamic and Scalable Tree (DST) [12].

2.3.1 Daa

In coding theory, error-correcting codes initially developed to correct errors on noisy communication channel. Coding theory provides a mechanism, which reduces data redundancy from a message and ensures transmission of message to recipient error free. Code-based cryptosystems such as McEliece cryptosystem [21] is the first cryptosystem developed in 1979, which has implementation based on binary Goppa code and its

security is based on difficulty of general decoding problem. Goppa codes forms a large family, providing a vast number of potential public keys, and there exists an efficient algorithm for decoding these codes in polynomial-time. Given a linear code [n, k] and encoded message, say vector 'y', it is hard to find another vector 'x', such that hamming distance between vector 'x' and vector 'y' remains minimal. Note that this corresponds to correcting a certain number of errors occurred on the code word 'x', represented by an error vector e, that is $y = x + e$. Using some secret parameter, McEliece cryptosystem transforms randomly selected binary Goppa code to a general linear code. It is hard to decode general linear code without knowledge of secrete parameters and it was proved to be NP-complete by Berlekamp, McEliece and van Tilborg in [23]. McEliece cryptosystem has been around as that of RSA cryptosystem, however it is almost not used in practice, due to large public key size of several hundred bits. Author K. W. Fan et al. [11] explains the power

2.3.2 DST

Author L. A. Villas et al. in [12] proposes a novel routing protocol called DST which adopts to any scenario. Different from static routing schemes, in DST, routes which are created during data transmission are not held fixed and also not dependant on the order of events. It aims to build a routing tree with the shortest routes in Euclidean distance that connects all source nodes to the sink node, maximizing data aggregation while reducing the distance connecting each coordinator node to sink. Routes are based on straight line segments, which are computed by coordinator nodes. Drawback of this protocol is that it does not explore the spatial correlation model. DST loses its performance in situations where nodes detect the same event.

III. Comparative Summary of Protocols

The following table shows the summary of characteristic of protocols of WSN.

Table I

Comparative Summary

Algorithm	Routing Structure	Objective	Data Aggregation Nodes	Overhead	Scalability	Spatial Correlation
Direct Diffusion	Tree Based	It constructs a spanning tree rooted at the Sink.	Clusterheads and Aggregatior Node	Medium	Low	No
SPT and CNS				High	Medium	No
LEACH	Cluster	It forms clustering of	Clusterheads and	Medium	Low	No

InFRA	Based	nodes and maximizes the overlapping of routes.	Intermediate Nodes	Very High	Low	No
DAARP				Medium	Medium	No
DDAARP				Low	Medium	No
YEAST				Very Low	Very High	Yes
DAA	Structureless	Efficient data aggregation without explicit maintenance of structures.	Clusterheads and Intermediate Nodes	Very Low	Very High	No
DST				Very Low	Very High	No

IV. Conclusions

Routing in WSN is challenging issue and has attracted a lot of researcher in recent years. In this paper we provide a survey of in-network data aggregation algorithms. In-network aggregation improves the lifetime of network by reducing the number of transmissions required for data collection which in turn reduces the energy consumption.

Author L. A. Villas et al. classify them into three categories, namely tree based, cluster based and structureless approaches. In this survey we have provided a basics of in-network aggregation and different protocols which fits in one of the routing structure.

Advantage of clustering structure over tree based and structureless approach is that they allow direct data aggregation at cluster head during data transmission phases.

References

[1] Chan, Haowen, Perrig, Adrian, and Song, Dawn, "Secure Hierarchical In-Network Aggregation in Sensor Networks," Department of Electrical and ComputerEngineering, Paper 12, 2006.

[2] NisheethShrivastava, ChiranjeebBuragohain, DivyakantAgrawal and SubhashSuri, "Medians and Beyond: New Aggregation Techniques for Sensor Networks," in Proceedings of the Second ACM Conference on Embedded Networked Sensor Systems, August 2004.

[3] E. Fasolo, M. Rossi, J. Widmer, and M. Zorzi, "In-network Aggregation Techniques for Wireless Sensor Network: A Survey," IEEE Wireless Comm., vol. 14, no. 2, pp 70-87, Apr. 2007.

[4] W. B. Heinzlaman, A. P. Chandrakasan, H. Balakrishnan, "Application specific protocol architecture for wireless microsensor networks," IEEE Transaction on Wireless Networking, 2002.

[5] L. Villas, A. Boukerche, R.B. de Araujo, A.A.F. Loureiro, "Highly dynamic routing protocol for data aggregation in sensor networks," Proceedings of the IEEE Symposium on Computers and Communications, ISCC10, IEEE Computer Society, Berlin (2010), pp. 496–502.

[6] S. Hougardy, H.J. Prömel, "A 1.598 approximation algorithm for the steiner problem in graphs," SODA99: Proceedings of the 10th Annual ACM-SIAM Symposium on Discrete Algorithms, Society for Industrial and Applied Mathematics, Philadelphia, PA, USA (1999), pp. 448–453.

[7] E.F. Nakamura, H.A.B.F. de Oliveira, L.F. Pontello, A.A.F. Loureiro, "On demand role assignment for event-detection in sensor networks," ISCC '06: Proceedings of the 11th IEEE Symposium on Computers and Communications, IEEE Computer Society, Washington, DC, USA (2006), pp. 941–947.

[8] L.A. Villas, A. Boukerche, R.B. Araujo, A.A. Loureiro, "A reliable and data aggregation aware routing protocol for wireless sensor networks," Proceedings of the 12th ACM International Conference on Modeling, Analysis and Simulation of Wireless and Mobile Systems, MSWiM09, ACM, New York, NY, USA (2009), pp. 245–252.

[9] K.-W. Fan, S. Liu, P. Sinha, "On the potential of structure-free data aggregation in sensor networks," in INFOCOM 2006, Proceedings of the 25th IEEE International Conference on Computer Communications, 2006, pp. 1–12.

[10] Leandro A. Villas, A. Boukerche, H. A. de Araujo and A. A. Luoureiro, "A Spatial Correlation Aware Algorithm to perform Efficient Data Correlation Aware Algorithm to Perform Efficient Data Collection in Wireless Sensor Network," Ad Hoc Network, 2011.

[11] S. Madden, M.J. Franklin, J.M. Hellerstein, W. Hong, "Tag: a tiny aggregation service for ad-hoc sensor networks," SIGOPS Operating System Review, 36 (SI), pp. 131–146, 2002.

[12] L.A. Villas, D.L. Guidoni, R.B. Araújo, A. Boukerche, A.A. Loureiro, "A scalable and dynamic data aggregation aware routing protocol for wireless sensor networks," Proceedings of the 13th ACM International Conference on Modeling, Analysis, and Simulation of Wireless and Mobile Systems, MSWIM10, ACM, New York, USA, pp. 110–117, 2010.