# **RESEARCH ARTICLE**

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# Man-In-The-Middle Attack Test-Bed in Vestigating Cyber-Security Vulner Abilities in Smart Grid Scada Systems

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# ABSTRACT

Theincreased complexity and interconnectivity of Supervisory Control and Data Acquisition (SCADA) systems in the Smart Grid has exposed them to a wide range of cyber-security issues, and there are a multitude of potential This presents SCADA-specificcyber-securitytestaccesspoints for cyber attackers. paper а bedwhichcontainsSCADAsoftware and communication infrastructure. This test-bed isused to investigate an Address Resolution Protocol (ARP)spoofing based man-in-the-middle attack. Finally, the paper proposes a future work planwhich focus esonapply ing intrusion detection and prevention technology to address cyber of the second secoer-securityissues inSCADA systems.

Keywords:SmartGrid,SCADA,Cyber-security,Man-in-the-middleattack

# I. INTRODUCTION

Supervisory Control and Data Acquisition (SCADA)

systemshavelongplayedasignificantroleinelectricali ndustry,becoming increasingly complex and interconnected as state-of-theartinformationandcommunicationtechnologiesarriv e.Theincreasedcomplexityandinterconnectionhasex posed them to a wide range of cyber-security vulnerablepoints.Inpractice,maliciousattackersordis gruntledemployees may gain unauthorized access to SCADA

systemsutilisingvulnerablepointsandthereafterlaunc helaborateattackswhichmayleadtocatastrophicdama ges.

The IEEE Standard 1402-2000 (R2008), 'Guide for ElectricPower Substation Physical and Electronic Security',

states:"Astheuseofcomputerequipmentwithinthesub stationenvironmentincreases, theneedforsecurity syst emstopreventelectronic intrusions may be comeeven m ore important."[1]

#### Inrecentyears, maliciouscyber-

securityincidentshavehappened from time to time. For example, in July 2010, theStuxnetwormattackedtheSiemensSIMATICWin CCSCADAsystem,usingatleastfourvulnerabilitiesof theMicrosoftWindowsoperatingsystem.Itisthemostf amousmaliciouscodeattacktohavedamagedanindustr ialinfrastructuredirectly [2].

In the early history of SCADA systems it was widely believed that such systems were secure since they were physically and electronically isolated from other networks. Stuxnetcrossed both the cyber and physical world by manipulating the controlsystemofthecriticalinfrastructure, demonstrat ingthat "security by obscurity" is no longer available pro ach.

With the development and deployment of SCADA systems,more and more cyber vulnerabilities will emerge in the SmartGrid. These vulnerabilities are not only from outside, such asterrorists, hackers, competitors or industrial espionages, butalso from utilities inside, such as ex-employees,

disgruntledemployees, vendorpersonnel for troublesh ooting, site engineers etc. In addition, cyber vulnerabilities in SCADAsystems result from well deliberate attacks as 28 inadvertentevents(e.g.,equipmentfailures,carelessne ss,andnaturaldisasters). Therefore, research on cyber-security issues for theuse of SCADA in the Smart Grid is extremely urgent and particularly significant as one of the keynote topics in thedevelopmentofsecure systems. However, researchinthis cross-

disciplinarysubjectisstillatanearlystage, and requires much more in depth investigation and analysis ofspecific vulnerabilities. To this end, the research presented inthis paper proposes a SCADA-specific cyber-security test-bedforsimulated cyberattacks. This environment provides a platform for the in-depth analysis of real attack scenarios, inorder to facilitate the development of effect iveattack countermeasure tools and technologies for the Smart Grid cyber domain.

Section II of this paper reviews he evolution of SCADAnetworks and their protocols in power

systems. Section IIIdescribes cyber-security vulnerabilities in SCADA systems and attack scenarios in a multilevel architecture. Section IVdiscusses related work about simulation, test-

bedandintrusiondetectiontechnologyforSCADAcyb er-security.Section V proposes a SCADA-specific cyber-security testbedth.etion.etion.com/

 $bed that investigates an Address Resolution Protocol (A\ RP) spoofing based man-in-the-$ 

middleattack.Finally,thediscussionandfutureresearc hworkarepresented.

# 1 Evolution of SCADA systems and protocol s

Tounderstandcyber-

securityissuesandchallengesinSCADA systems of Smart Grids, it is better way to brieflyreviewtheevolutionofSCADAsystems.From1 960stotoday,theSCADAsystemshaveundergonethre emainphasesofdevelopment,i.e.,central,distributeda ndnetworkedarchitecture [3].

# Centralarchitecture

ThisisthefirstgenerationSCADAarchitectur einwhichmainframe systems with redundancy are in charge of all thefunctions such as Remote Terminal Unit (RTU) polling, dataprocessing,display,report,dataarchiving,andrun ningapplicationprograms.ThecommunicationofSC ADAisrealizedbyvendor-

proprietaryequipmentandprotocol.

#### Distributedarchitecture

Startinginthe1980s,themajorityofSCADAs ystemsadopta distributed architecture in which multiple computers in anetwork(e.g.,LocalAreaNetwork(LAN))shoulderth ecomputing burden together and different computers

realizedspecificfunctionsandroles.Comparingwithth efirstgeneration SCADA systems, the communication protocols aresimilar. Hence, the systems are still limited by the different/vendors.

#### Networkedarchitecture

ThethirdgenerationSCADAarchitectureiso pensystemarchitecture, rather than a vendor proprietary environment, which utilizes open standards and protocols and distributesSCADA functionality across Wide Area Network (WAN) andnotjustLAN.Comparing with the first and secondar chitectures, the major improvement lies in the application

ofWANprotocols(e.g.,theInternetProtocol(IP))forco mmunication.

WiththedevelopmentofSCADAarchitectur e,communication protocols in SCADA system have also beendevelopedfrompoint-topointlinktoopenandstandardprotocols.Fig.1illustrate s

abriefsurveyoftheSCADAprotocoldevelopmentfro m1970s.



Fig.1:Thedevelopmentofcommunicationprotocolsin

proprietaryhardware,softwareandcommunicationspr otocols.However,theinteroperability,connectivity,a ndcompatibilityofmodernSCADAsystemsbringhug echallenges in order to make the current systems more secure from unintentional or malicious cyber attacks. In addition, since the lifecycle of SCADA equipment is 15-20 years, it is not uncommon that 'smart' and 'dumb' devices coexist in thefield. Therefore, it is of importance to understand thei ntegrationissuesbecausebothpastandfutureSCADAp rotocolsarecombined in theSmartGrid.

# 2 Cybervulnerabilitiesandattackscenarios inSmartGrid SCADAsystems

# Cyber-securityinSCADAvs.ITsecurity

In current industrial and academic fields in terms of cyber-

securityofcontrolsystems(e.g.,SCADA),powersyste mresearchers may not master the knowledge which IT

securityexpertsknow, and vice versa. Infact, there arem any differences between the two areas which will depen dondifferent countermeasures for cyber vulnerabilities. Table 1 describes the comparison of SCADA cyber-security and IT security [4].

Subject	SCADACyber-	ITSecurity	
	security	-	
Availability	Veryhigh	Lowto moderate	
Integrity			
Confidentiality	Low	High	
Authentication	High	Moderate	
Applicationof	Sloworeven	Frequent	
Patching	impossible		
Anti-virus	Uncommon	Commonly	
TechnologyLifetime	15-20years	3-5years	
Timecriticality	Critical	Delaystolerated	
Communicationproto	IEC 61850,		
cols	IEC60870-	TCP/IP,UDP	
	5/6,DNP3,		
	Modbusetc		
Computingresources	Verylimited	Unlimited	
Cyberforensics	Limited, if any	Available	
Securityawareness	Poor	Good	
Impactsofsecurityco	Economicimpacts,		
mpromise	equipment	Economicimpacts	
	damageandperson		
	nelsafety		

Table1:ComparisonofSCADAcyber-securityandITsecurity

### Cybervul nerabilities and consequences

AccordingtothereportoftheUSNationalInstituteofSta ndards and Technology (NIST) [5], there are three maincyber-security requirements for SCADA systems in the SmartGrid:availability,integrity and confidentiality.

Anintentionalviolationofacybersecurityrequirementis SCADAsystems FromthehistoryofSCADAsystems, it is interred that S CADA systems from the 1960s to 1980s could probably

besecurefromcyberattacksbecausethesystemsutilize calledanattack.Sometypicalcyberattackswhichmayc ompromise SCADA systems in the Smart Grid are listed in[6], such as denial-of-service (DoS) / distributed DoS (DDoS),malicioussoftware,identityspoofing,passwo rdpilfering,eavesdropping,intrusion,sidechannelattacks.Table2

demonstratespossiblecyberattacksandconsequencesi nSCADAsystems.

Cyberattacks		Consequen ces			
·		Cyberspace	SCAD A		
DoS/DDo	Crashservic es	Comprom iseavailabi litvand	Disable the monitoringandcont		
S	Floodservic es	integrity;u nresponsi ve nodes	rolsystem;lossof load; loss ofinformation.		
Intellige ntattac ks	attackpro tectiverel aysetting	Comprom iseintegrit y	Trigger cascadingeffectswh ichmayresultinama jorpoweroutage thatcanbecatastrophi c.		

Intrusion	IPscans Portscans	Compromis econfidenti alityandinte grity	Control aspects of thebehaviour of the systemat intruders' will whichmayleadtolos sofload.
Malicio ussoft	Virus Worms Trojanhorse S	Compromis eavailabilit y,integrity	SCADA systems arecompromised e.g., slowdown thecommunication
ware	s Backdoors	orconfident iality	betweensubstations andcontrol centres.
Identit ySpoof ing	man-in-the- middle message replays network spoofing	Compromis econfidentia lity,integrit y	Causesafetyissue sinSCADA systems byimpersonating anauthorizeduser
Passwo rdPilfe ring	sociale ngineerin g	Compromis econfidenti alityor accesscontr ol	The severity ofconsequences dependsonthelevel ofauthority intermsofth enassword

Table2:CyberattacksandconsequencesinSCADAsystems

#### **CyberAttackScenarios**

We propose that most cyber attack scenarios usually fall intothefollowingfivelevelarchitecture(LeverI-V).Severalattack scenarios which may cover more than one level aredescribedtoillustratehowinsidersandoutsiderscou ldexploitthe vulnerabilitieslisted inthe table 2.

#### 1) LevelI:Thetarget

rangeofcyberattacksinthemostessentiallevelofthefiv e-levelarchitecturecontainsasubstation and field devices, such as Intelligent ElectronicDevices (IEDs), Human Machine Interface (HMI), LAN etc.Theattackpathsinthislevelcontainwiredchannels andwirelesschannels.Forexample,unauthorizedusers orattackers may access to the LAN in a substation by dial-up,Virtual Private Network (VPN) or wireless and then launchmaliciousattackssuchasportscan,sniffing,man -in-the-middleattack etc.

Scenario 1: Simulated attacks on an IEC 61850 based IED inan experimental setup is presented in [7], such as DoS attack,passwordcrackattackandARPspoofingattack. Aftersuccessful cyber attacks, an attacker can access the SubstationConfiguration Description (SCD) file including the electricdiagramofthesubstation,communicationinfra structure, configure of IED and IEC 61850 based traffic. Therefore, the

attackermaylaunchmaliciousactionstooperatecircuit breakersbasedon identifiedinformation[8].

Scenario 2: Disgruntled employees or other attackers may beable to launch man-in-the-middle attack to sniff and interceptnetwork traffic between master station and slave station basedonMODBUSprotocolintheLANwithinasubstat ion.Moreover,attackerscansendfalseinformationtoth eMODBUS master or slave by poisoning the MODBUS

packetaddresses.Itcanbeutilizedtorealisefalseread/w ritecommands to the MODBUS server, block the communicationbetween the server and the client, restart the server, or evenshutdownpartofthe gridetc [9].

2) Level II: The scope of attack targets in level II covers the control centre and the communication between the substationand the control centre. The widely used com munication protocols are Distributed Network Protocol Version 3 (DNP3) and IEC 60870-5 serials which are lack of mature security mechanism.

Scenario3:Anunscrupulousattackermayabletouseide ntify spoofing attacks to obtain network traffic in

level

II.Forexample,captureddatareflectingnormaloperati onsinthecontrolcentreisplayedbacktotheoperator.Itl eadstothe operator's HMI to appear normal and consequently theattack will not be recognized. In addition, the attacker couldcontinue to send malicious commands from the control centretofielddevicesinlevelI,whichmaycauseundesir abledamageswhiletheoperatorremainsunawareofthe realsituationof theSCADA system[9].

3) Level III: This level mainly focuses on communicationbetweenutilitycontrolcentreswhich maybelongtotransmissionordistributionoperators, In dependentSystemOperators(ISO)orpowerplants. Th emostpopularandexposed SCADA protocol in this level is Inter-control CentreCommunication Protocol (ICCP).

Scenario 4: According to the LiveData ICCP Server

whitepaper[10],LiveDataICCPServercontainsakind ofvulnerability, i.e., heap-based buffer overflow. The

LiveDataimplementationofRequestforComments(R FC)1006isvulnerabletoaheap-based

bufferoverflow.Bysendingaparticularlycraftedpacke ttoavulnerableLiveDataRFC1006 implementation, an attacker may trigger the overflow toexecute malicious code or crash a LiveData ICCP Server tocauseaDoSattack.

4) Level IV: This level covers Demilitarized Zone (DMZ), anetworksegmentasa"securitybufferarea",connectin

gcontrolcentreswith corporatenetworks.

Scenario5:Itispossibletoestablishcommunicationcon nection between corporate networks and control centres,forinstance,communicationbyTCPacknowle dgementpackets.Furthermore,communicationpathss upportedbyvendors may be utilized by attackers to launch cyber attacksfromthecorporationnetworktothecontrolcentr e.After

accessingtothecontrolcentre, attackersmaygainanyco nfidentialandessentialinformation, modifycontrolco mmands, and tamper the configuration files.

5) Level V: This level is the outermost layer in the multilevelarchitecture which includes corporate networks and connectedInternet. The cyber vulnerabilities and attack scenarios of thislevelalmost belong to ITnetwork securityarea.

Scenario 6: A utility employee who can access to computerinformation service may install or run a

orseemingly computer "game" innocuous application software from a friend, exemployee, vendoror actually any one with legitimate co nnection to the employee's utility. The installed softwareincludes a Trojan horse program which opens a backdoor into he computer network. The attacker who invents the Trojanhorseprogramcangainaccesstothecomputerint hecorporate networkfrom Internet.andfurtherlaunchmanykindsofattackssucha sDoS.Thecomputerinformationsystemsin thecorporatenetworkarenowinjeopardy[11].

# Relatedwork

Increasingly, academic and industrial related organisationsare focusing on cyber-security issues of SCADA systems intheSmartGrid.However,crossdisciplinaryresearchconnecting developments in power systems and IT still hassomewaytogo.Inthissection,relevantpublishedlit eratureintermsofsimulation,testbedandintrusiondetectiontechnology for SCADA cyber-security research is surveyedandsummarized.

#### Simulationandtest-bed

It is indispensable to set up a simulation platform or test-bedfor research on SCADA cybersecurity in the Smart Grid, especially for cyber vulnerability and risk assessment, and interaction and interdependence between power system and cyber infrastructure [6]. Moreover, there is a lack of practical, statistical and historical data about cyber-security toward theelectricinfrastructure.

Oneapproachtoobtainpracticaldataistobuildacompar ativelysimplesimulationwhichcanapproximateareals ituation,forexampletheUSIdahoNationalLabSCAD A test-bed [12]. The paper [13] also introduces a test-bedforSCADAcyber-

securityinwhichtheexperimentillustrates the vulnerability of the network client to a DDoSattackandtheabilityoffilteringto mitigate an attack.

Inaddition,theEuropean6<sup>th</sup>FrameworkProgram(FP6) project<sup>c</sup>CriticalUtilityInfrastructureResilience<sup>2</sup>(CR UTIAL)

[14] set up two test-beds for tele-control and microgrid tocollect data statistics and evaluate malicious attacks in gridtele-operationandmicrogridcontrolscenarios.

Furthermore, researchers from the University of Arizona inUS[9]developedatestbedtoanalysethesecurityofSCADAcontrolsystems(T ASSCA).Thetest-bedadopteda

TCP, Modbus and DNP3 protocol analyser to

#### detect

SCADAattackanomalies, for example protocol statetr ansition analysis.

Other researchers have tried to exploit the coupled power

gridcommunicationnetworksimulatorbasedonsoftw areagentsor application program interface (API) methods [15, 16] usingcommercial-off-theshelf(COTS)simulationtools,suchasMATLAB,

#### PSCAD/EMTDC,

OpenDSS,PSSTMNETOMAC,NS2/3,OP NET,OMNET++etc.

From published work and the above examples it is known thatauthentic simulation and accurate testbeds are effective toolsforSCADAcybersecurityresearch.However,comprehensive and welldeveloped tools require significanteffort to fully develop but are often propriety, hence limitedopen simulation and test-bed resources are available to thewiderresearch community.

#### Intrusiondetectiontechnology

Before full deployment and operation, SCADA systems in theSmart Grid will inevitably contain legacy systems that cannotbeupdated,patched,orprotectedbymanytraditi onalITsecuritytechniques.Withlimitedcomputingres ourcesinlegacydevicesandevennosecuritydesignfor SCADA systems, it is difficult to embed traditional cybe rsecuritytechniques into the Smart Grid with legacy systems. In these situations, a feasible approach is intrusion deploy to detectionandpreventionsystemfor SCADAsysteminSmartGrid.

Intrusion detection technology in the IT domain is relativelymature.Numerousintrusiondetectionmetho dshavebeenpresented[18]andsomeofthemhavebeena ppliedintoSCADA systems [19, 20]. However, research on this cross-disciplinesubjectisstill atanearly stage.

Theprimarylimitationofthecurrentintrusion detectionsystems for SCADA is a lack of adequate knowledge and experience of SCADA applications and protocols. The USIdaho National Laboratory indicates [19] the above limitationintermsofcurrentIntrusionDetectionSyste m(IDS)application to SCADA systems, and then presents the futureSCADA IDS technologies implementing signature matching, flow analysis, and data inconsistency detecti ontailoredparticularly for SCADA systems. However, there is lack of experimental study.

A. Carcano et al. proposes a state-based intrusion detectionsystem for SCADA system based on Modbus/DNP3 protocols[21,22].ThepresentedIDScontainsbothtrad itionalsignature-basedtechniquesandanovelstateanalysistechniqueswhichcanmonitorcriticalstatesan didentifycomplexcyber attacks.

The model-based detection is not new in traditional IDS, e.g., specification-based intrusion detection can be seen as modelbased. [23] believes that model-

basedmonitoringfordetectingunknownattacksismore feasibleforSCADA systems than general enterprise networks. The paper describesthree model-based techniques for monitoring Modbus TCPnetworks, i.e., protocol-level modes, communication-patternbaseddetectionandlearning-basedapproach.

#### Arule-

basedintrusiondetectionsystemforanintelligentelectr onic device (IED) based on IEC 61850 is realised bySnort in [7] which develops rules by using experimental databased upon simulated cyber attacks, such as denial of service(DoS) attack, password crack attack and ARP spoofing attack.However,therulesdonotrefertoIEC61850prot ocolanalyses. In addition, several other new approaches have beenpresented to deal with intrusion and anomaly detection, suchas a neural network based [24] and rough sets classificationalgorithm[20].

# **3** SCADA-specificcyber-securitytestbedandsimulatedattacks

In this section, a clear and effective SCADA-specific cyber-security test-bed focusing on the core level of the five-levelarchitecture in presented Section III is in order to investigatecyber-security vulnerabilities in SCADA systems. The test-bedisbasedonarealgridconnectedphotovoltaic(PV)SCADAsystemthathasb eenimplementedinapracticalenvironment.

#### **Test-bedarchitecture**

In Fig. 2, the test-bed architecture contains five computers (A-E)andaswitch.Threewindows-basedhosts(A,B,C)simulatereal-

timeSCADAcommunicationinarealsubstation LAN. The host A simulates the master station orHMI where COTS SCADA supervisory control software isinstalled.ThehostBwithanotherCOTScommunicati onsoftware is a protocol gateway. The two hosts A and B areconnected by a switch. IED simulator communicates with

theprotocolgatewaybyIEC60870-5-

103protocol.Duetoconfidentialityandsecurityconcer nsthenamesoftheSCADA software, the switch and

The

the simulated IED in thetest-bedare withheld.

The Linux-based host D is utilized to simulate an intrudedcomputer inside the LAN or any possible laptop connected tothe LAN from the outside (e.g., maintenance laptop access),which can be illegally control by an attacker. Many cyberattacks can be investigated in the test-bed such as DoS attack,ARPspoofingattack andman-inthe-middleattack.

In addition, a SCADA-specific IDS based on The InternetTraffic and Content Analysis (ITACA) will be realised in theLinux-based host E which is connected to the LAN by portmirroring.ITACA[25]isasoftwareplatformfortra fficsniffer and real-time analysis of IP network which has

beendevelopedbyCentreforSecureInformationTech nologies(CSIT) in the Queen's University of

#### Belfast.

extendableanalysistoolenablestheimplementationof plug-intoperform specifictask,e.g.,IDS.Inthe testbed,SCADA-specific IDS will be created using a well defined C/C++ APIbasedonITACAplatform.

#### Vulnerabilityscan

A comprehensive vulnerability scanning program, Nessus, isused to scan potential vulnerabilities by host on the testedsystem. The master host, the protocol gateway host and theswitch are scanned by Nessus [26] and the scan results

arelistedinTable3.Forexample,boththemasterandthe protocol gateway have critical vulnerabilities, i.e., MicrosoftWindowsServerMessageBlock(SMB)vul nerabilities,which may allow an attacker to execute arbitrary code orperformadenialofserviceagainstthe remotehost.



Fig.2:SCADA-specificcyber-securitytest-bedarchitecture

Hos t	Severity	Descripti on
MasterW	Critic al(10 .0)	Microsoft Windows ServerService Crafted RPC Request HandlingRemoteCodeExecu tion
A(193.100.10 0.98)	Critic al(10 .0)	Microsoft Windows SMBVulnerabi litiesRemoteCode Execution
	Medium(5. 0)	MicrosoftWindowsSMBNU LL SessionAuthentication
	$\begin{array}{c} \text{Medium}(5. \\ 0) \end{array}$	SMBSigningDisabled

<b>D</b> 1	Critic al(10 .0)	Microsoft Windows Server ServiceCraftedRPCReques tHandlingRemoteCodeExe cution
GatewayWin dows host B(193,100,10	Critic al(10 .0)	Microsoft Windows SMB VulnerabilitiesRemote CodeExecution
0.80)	Medium(5. 0)	MicrosoftWindowsSMBNU LL SessionAuthentication
	Medium(5. 0)	SMBSigningDisabled
	Medium(6. 4)	SSL Certificate Cannot Be Trusted
Switch(19	Medium(6. 4)	SSLSelf-SignedCertificate
3.100.100.254 )	Medium(5. 0)	SSLCertificateExpiry
	Medium(4. 0)	SSLCertificate Signed using WeakHashingAlgorithm
	Low(2.6)	SSL / TLS Renegotiation HandshakesMiTMPlaintext DataInjection

Table3:Cybervulnerabilitiesbyhostinthetest-bed

#### Man-in-the-middleattackusingARPspoofing

The ARP is primarily used for resolution of network layeraddresses(IPaddresses)intodatalinklayeraddres ses(Ethernet Medium Access Control (MAC) addresses) in LANcommunication. In order to obtain the MAC of a destinationhost, a source host broadcasts an ARP request to all hosts inthe LAN asking for the MAC address of the destination withIP address. The destination host responds with given IP andMACinanARPreply.Thesourcehostcachesthe<I P,MAC>pairing in local ARP cache table so that it

P,MAC>pairing in local ARP cache table so that it does notneedbroadcastthe same request in thenear future.

The ARP spoofing attack is used to modify the cached <IP,MAC>pairing in the local ARP cache table. An attacker canassociate a malicious host's MAC address with IP of a targethost by modifying its ARP cache to add/update an entry withan <IP, MAC>mapping, so that the attacker can launch DoSattack, perform man-in-the-middle attack and gain access toconfidentialinformation[17].

The man-in-the-middle attack allows an attacker to sniff aLANbyARPspoofing.Firstly,theattackerredirectsc ommunicationtrafficbetweentwovictimhoststothem alicioushost.Then,themalicioushostwillsendtherecei

ved or modified packets to the original destination, sothat the communication between the two victim hosts

looksnormalandthevictimsmaynotnoticethattheirco mmunication information has being sniffed by the attacker[30]. Actually, Stuxnetcan also be described as a man-in-the-middleattackwhichfeeds monitoring software fakeinputreadings.

In the test-bed environment presented in this paper, an ARPspoofing attack is launched by a Metasploit[27] module inBacktrack 5 [28] which is Linux-based penetration testingsoftware.Alsoexaminedare,Cain,Ettercap,Ser ingeetc,whichareeffective toolsto launchARP poisoningattacks.

Normally, when the communication protocol gateway (IP:193.100.100.80, MAC:\*\*:\*\*:43:bb:74:4a) wants to send information from the slave station, such as remote measurevalues, remote communication values, or collection of electric energy, to the master station (IP:193.100.100.9 8, MAC:

\*\*:\*\*:43:b7:b9:90), it broadcasts an ARP request in the LAN"Who has 193.100.100.98? Tell MAC: \*\*:\*\*:43:bb:74:4a".All the other hosts in the LAN receive the request. However,only the host Aanswers back in an ARP reply "I have IP193.100.100.98, My MAC is \*\*:\*\*:43:b7:b9:90". And thenthehostB updatelocal APRtable using the <193.100.100.98, \*\*:\*\*:43:b7:b9:90>mapping.

However, ARP is a stateless and trusting proto coland does not provide any verification mechanism to v erify the authenticity of the ARP requests and replies, so ARP attacks are possible launched by malicious hosts in a LAN. In the ARP cache poisoning attack launched by Metasploit , the attacker (host D) sends ARP replies to the protocol gateway (host B) indicating that the materstati on (host A) with the IP

193.100.100.98 MAC has the \*\*:\*\*:27:ed:09:0fwhich is theMAC address of the so the host B will update attacker, itsARPcachetablewiththe<193.100.100.98, \*\*:\*\*:27:ed:09:0f>paring. The results of the APR spoofingattack are recorded in Fig. 3 by Wireshark[29]. 3 Fig. showsthattheattacker(hostD)impersonatesthemaster station(hostA)sothattheprotocolgateway(hostB)will sendpackets destined to the master station to the attackerinstead.

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No. Time		Source	Destination	Protocol	Info					1. C
899 180,83	2724	Cadmisco_edi09:0f	Del1_bb17414a	ARP	103.10	0,100.9	8.15	101	08:00127:edi0	1010
901 180. 93	6915	CadmusCoed:09:0F	Dell_bb:74:4a	ARP	193.10	0,100.9	8 15	at.	08:00:27:ed:0	9:01
903 181.03	0032	Cadmusco_ed:09:0f	Dell_bb:74:4a	ARP	193.10	0.100.9	6.15	78	Off:00:27:ed:0	9:07
907 183.7	3738	cadausco_ed:09:0f	:bell_bb:Fd:da	ARM	193.10	0.100.9	8.14	at	08:00:27:ed:0	9:0f
416 185.00	10111	CARMINESCO., ed. UV . OF	Dell's hits 74-44	ARP	191.10	10.100. *	8 11	11	0.01001221.0110	1101
918 185,14	10232	CadmusCo.,ed:09:0f	0e11_bb:74:4a	ARP	193.10	00,100.9	4.15	HL.	08:00:27:ed:0	2:0f
920 185.53	4429	CadnusCo_ed:09:0f	Dell_bb:74:4a	ARP	193.10	0.100.9	5.14	at	08:00:27:ed:0	9:0f
924 185, 6	13942	Cadmusco_ed:09:0F	De11_bb17414A	ARP	193, 10	00,100,9	8 15	48	081001271#810	9101
930 186.11	2925	CadmusCo_ed:09:0f	Dell_bb:74:4a	ARP	193,10	0.100.9	1 15	AL.	08:001271##10	010f
933 188.74	\$5559	Cadmusta_ed:00:0f	Dwll_bb:74:4s	ARP	192.10	0.100.9	·清.清末	at	OB:00:27:#d:0	2:0f
		1								
T. Colones Dite	1.60	bates on when LAND	NAME AND ADDRESS OF	botuned	7480 h	Sec. 1				

Fig. 3:The results of ARP spoofing attack on the protocolgateway(hostB) inWireshark

Similarly, the master station (host A) can also become thetarget host of ARP spoofing attack. After local ARP cache inthe master is poisoned, the <IP, MAC>pairing in the ARPcache table will be updated from <193.100.100.80, \*\*:\*\*:43:bb:74:4a> to

<193.100.100.80.

\*\*:\*\*:27:ed:09:0f>.TheresultcanalsobeseenfromFig .4. Furthermore, by poisoning the master station (host A) and theprotocol gateway (host B) at the same time, the attacker (hostD) can silently stay in the middle of the two hosts to launchman-in-the-middle attack in the test-bed so that the attackercan easily sniff all the traffic sent in both directions and injectnewdataintoboth.

🛃 ARP attack - Weesh	ark			0.0	. Long
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No. Time	Source	Destination	Protocol	lefe .	-
1377 218.801116 1383 219.243129 1386 219.346599	Cadmisco_ed:09:0F Cadmisco_ed:09:0F Cadmisco_ed:09:0F	Dell_h7:b9:90 Dell_b7:b9:90 Dell_b7:b9:90	ARP ARP ARP	193,100,100,60 is at 08:00:27:ed:00:0f 193,100,100,80 is at 08:00:27:ed:09:0f 193,100,100,80 is at 08:00:27:ed:09:0f	
1392 219,457019	cadeusco.ed10910f	De11_b7109190 De11_b7109190	ARP	193,100,100,80 is at 081001271ed10910f 193,100,100,80 is at 081001271ed10910f	
1402 221,190175 1407 222,316476 1414 222,524055	CadmusCo_ed:09:0F CadmusCo_ed:09:0F CadmusCo_ed:09:0f	Dell_07:09:90 Dell_07:09:90 Dell_07:09:90	ARP ARP ARP	193.100.100.80 is at 08:00:27:ed:09:0f 193.100.100.80 is at 08:00:27:ed:09:0f 193.100.100.80 is at 08:00:27:ed:09:0f	
1423 223,413771 1425 223,518366	cadmusco_ed:09:0f Cadmusco_ed:09:0f	bel1_b7:b9:90 bel1_b7:b9:90	ARP ARP	193.100.100.80 1s at 08:00:27:ed:09:0f 193.100.100.60 1s at 06:00:27:ed:09:0f	
Frame 1300; 6 Ethernet II, Address Resol	0 bytes on wire (48 srci cadmusco_edi09 ution Protocol (rep	0 bits), 60 byt 10f ((****)2716 19)	es capt d:09:0f	ured (480 bits) ), Osti Dell_b7:b9:90 (1**#* 43:b7:b9:90)	,

Fig.4:TheresultsofARPspoofingattackonthemaster(host

#### A)inWireshark

Fig. 5 and Fig. 6 illustrate that the attacker can easily obtain the communication information between the master and theprotocol gateway in the test-bed. For examples, frame 3499 inFig.5describesachangedremotemeasurevalue from theIED to the master, and frame 1967 in Fig. 6 shows a remoteoperation command from the master to the protocol gatewayin the SCADA system. The malicious attacker may utilize theinterceptedinformationtolaunchmoresevereattac kslater.

Man in the middle attack - Wresh	44		10 (B
Eds fids You fo Supture &	pratyce Statistics Talephony Jools	Helb	
	*	2 00 QQQ	E 👹 KA 🥵 🌾 😫
Filter	)	* Expression One Apply	
tan.         Terme         Source           3444         408.219290         109.100.1           3494         609.116212         109.100.1           3494         609.255489         189.100.1           3499         609.116212         109.100.1           3499         609.851847         193.100.1           3500         10.854889         189.100.1           3501         10.854831         193.100.1           3502         10.854831         193.100.1           3501         10.854831         193.100.1           3517         411.9674831         193.100.1           319         82.054889         193.100.1	Destination         Path           100.80         193.100.100.94         004           000.80         193.100.100.94         006           100.80         193.100.100.94         006           100.80         193.100.100.94         006           100.80         193.100.100.94         006           100.80         193.100.100.94         006           100.80         193.100.100.94         006           100.80         193.100.100.94         006           100.80         193.100.100.94         006           100.80         193.100.100.94         006           100.80         193.100.100.94         006           100.80         193.100.100.94         007           100.80         193.100.100.94         007           100.80         193.100.100.94         007           100.80         193.100.100.94         007	cel bis Source port: 17682 Source port: 17682	Destination port: 17682 Destination port: 17682
Frame 3499: 70 bytes on Ethernet II, Src: Dell Internet Protocol, Src: User Satagram Protocol, Date (28 bytes)	wire (360 bits), 70 bytes Mo:74:44 (++++ 193.100.100.40 (19).100.100 SFC Port: 17682 (17682), D	captured (560 bits) ), Dit: cadeusco_ed:00: 0.80), Dit: 193.100.100 St Port: 17882 (17882)	0f ()

Fig. 5: The sniffed traffic from the protocol gateway (host B)to the master (host A) captured by Wiresharkin the attacker(hostD)

iter			· Expression Clear Apply		
n. Toter 794 282.466404 849 290.899079	Source 193.100.100.98 193.100.100.98	Destination Protoc 193,100,100,80 UDP 193,100,100,80 UDP	source port: 17682 Source port: 17682	Destination port: 17682 Destination port: 17682	1
885 293 947786	193.100.100.98	193,100,100,80 UD#	Source port: 17682	Destination port: 17682	
951 306 043121	193.100.100.98	193,100,100,80 UD#	Source port: 17682	Destination port: 17682	
967 309 356031	193.100.100.98	193,100,100,80 UD#	Source port: 17682	Destination port: 17682	
55 111.097631	191,100,100,96	193,100,100,80 LDP	Source port: 17682	Destination port: 17682	
542 321.190727	193,100,100,96	193,100,100,80 LDP	Source port: 17682	Destination port: 17682	
547 322.738311	197,100,100,96	193,100,100,80 LDP	Source port: 17682	Destination port: 17682	
122 532, 133806	193,100,100,98	193.100.100.80 UDP	Source port: 17682	Destination port: 17682	
175 341, 422139	193,100,100,98	193.100.100.60 UDP	Source port: 17682	Destination port: 17682	

Fig. 6: The sniffed traffic from the master (host A) to theprotocolgateway(hostB)capturedbyWiresharkintheattacker (hostD)

# II. DISCUSSION AND FUTUREWORK

Intheman-in-the-

middleattackexperiment, anattack simulator is developed by C/C++programming which can send modified information such as remote op erationcommands, remote measure values and remote c ommunication values to the master station or the protocolgateway. The injected malicious data from the attacker willdisplay on the screen of the master host which mav misleadtheoperator'sjudgement.Evenworse,thefalse remoteoperation command such as "open the circuit breaker" from he attacker will make the PV grid lose loads and decreasepowersupplyreliability, and eventhreat perso nalsafety.

The above ARP spoofing based man-inthe-middle attack inthe test-bed belongs to the level I of five-level architecture inSection III. The indicator of the presence of ARP poisoningbased man-in-the-middle attack can be detected by indepthpacket analysis (e.g., IDS), because packets sniffed by theattacker have unmatched the <IP, MAC>pairing. Actually,SCADA-specific IDS can address not only known man-in-the-middleattackinthetest-

bedbutalsounknowncyberattacks.

According to the aforementioned survey, investigation

and discussion, authors' next research planistodevelop a SCADA-

specificIDSagainstbothknowncyberattackssuch

as man-in-the-middle attack and novel cyber attacks in theSmartGrid environment.

# **III. CONCLUSION**

According to the evolution of SCADA systems and cybervulnerabilitiesandattackscenariosinpublishedli terature, it is clear that a large number of potential cyber-security issuesare increasingly probable on SCADA systems in the SmartGrid. This paper provides overview of cyberan the securityvulnerabilities of SCADA systems in Smart Grid. The paperhas also proposes a SCADAspecific cyber-security test-bedthat investigates an ARP based man-in-thepoisoning middleattack.Fromtheexperimentresults,itisinferred

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that malicious cyberattacks such as man-in-the-

middlecaninfluence and compromise secure and reliable operation of SCADA systems. Therefore, research on cyber-

securityvulnerabilities for SCADA in the Smart Grid is

extremelyurgentandparticularlysignificantasoneofth ekeynotetopicsinthedevelopmentofsecuresystems.F inally,thepaperpresentsthatSCADA-

specificIDS is a promising approach to address cyber vulnerabilities in SCADA systems in authors' future research work.

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