# **RESEARCH ARTICLE**

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# **The Skien Hash Function**

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# ABSTRACT

In this paper, we simulate the performance of skein hash function in VHDL FPGA. Skein is a new cryptographic family. It is designed in such a way that it is simple, easy and flexible. This function which is implemented in software based. This is a hashing function which provides security.

**Keywords:** Skein hash function, Threefish block cipher, MIX function, Key schedule, unique block iteration (UBI).

#### I. Introduction

The new family of cryptographic hash function is skein. The advantages of this hash function are speed, security, simplicity and flexibility due to its design as well as easy to analyse. Skein hash function is fast and secure Skein has three different internal state sizes as 256, 512, 1024 bits and output may be any size. It is a good replacement of SHA family. Skein is an efficient tool to be used for large number of functions due to its optional and extendable argument system [1]. Skein is efficient for both hardware and software platforms.The remainder of this paper is organized as follows:

Section II presents the Basic structure of skien, Section III discusses about the skein algorithm, and Section IV draws the conclusion.

## II. Basic structure of Skein

#### 2.1. Overview of Skein

Skein is a new family hash function with three internal state sizes: 1024,512,256 bits.

- Skein- 512 which is used for all present hashing applications with secures and remains secure for the future.
- Skien-1024 is ultra conservative variant [1]. Skein-1024 is the twice of skein- 512. Skein 1024 is twice as fast as Skein- 512.
- Skien-256 is low memory variant [1]. The above discussed internal state size which can support any size of output. Replacement of one type of hash function with other type of hash function shown in Table: 1[1]

		State	Output
Replace	With	Size	Size
MD5	Skein-256-128	256	128
	Skein-512-128	512	128
SHA-1	Skein-256-160	256	160
	Skein-512-160	512	160
SHA-224	Skein-256-224	256	224
	Skein-512-224	512	224
SHA-256	Skein-256-256	256	256
	Skein-512-256	512	256
SHA-384	Skein-512-384	512	384
	Skein-1024-384	1024	384
SHA-512	Skein-512-512	512	512
	Skein-1024-512	1024	512

# Table: 1: Replacements for MD5, SHA-1 and SHA-2

#### 2.2. Basic building blocks

Three new components are needed to build Skien.They are three fish block cipher, Unique block iteration (UBI), Optional argument system.

# > Three fish Block cipher

It is a tweakable block cipher at the Skein's Core with any internal size as 256, 512, 1024 bits.

#### > UBI- Unique Block iteration

This is the type of chaining mode which uses three fish in order to build the compression function which maps any input size to any or fixed output size.

# > Optional argument system

This makes skein which supports any optional Features.

#### 2.2.1. Threefish Block cipher

Threefish is a large tweakable block cipher [66]. It is also defined for three different block sizes as 256, 512, 1024 bits. The key which should be same size as the block and the tweak value is only 128 bits which is applicable for all the three different block sizes.

The main principle of Threefish is that more number of simple rounds which are secured than few numbers of complex rounds. There are three different mathematical operations used by Threefish block cipher are Exclusive-OR (XOR), Addition and constant rotation.



Fig: 1. MIX Function

The function in three fish which is actually a MIX function. **Fig: 1[1].** Each function which consists of a single addition, a rotation by one constant, an Exclusive OR function.



Fig: 2 .Three fish-512.

Figure.2 which shows the operation of MUX function to build threefish-512.Each skien-512 has 72 rounds which has four mix functions which gets followed by a permutation of eight 64 bit words. Each function made on 64 bit words. At every four rounds one sub key will be generated. The word permutation is same for every round. The rotations here are constant which are chosen in such a way to maximize diffusion [1].

These rotations are constant for every eight rounds. The sub keys are generated from the key schedule. The key schedule generates the sub keys by use of key and tweak. Each sub key contributes key words, tweak words and counter value. To create a key schedule, key and tweak words which are extended with one parity word extra which is XOR of all the words .For threefish-256 there are 72 rounds and 2 MIX functions followed by permutation.

For threefish-1024 the rounds are 80 and have 8 MIX functions. The rotation constants and permutations are different for different internal Block sizes. But it should maximize the diffusion.



Fig: 3.Threefish-256

# 2.2.2. Unique Block Iteration (UBI)

It is a chaining mode which produces fixed output by using input chaining value and arbitrary length input string. The tweak is the heart of UBI .Every block is processed with unique variant of the compression function. UBI chaining mode which ensures the above work. UBI which follows Matyas-Meyer-Oseas [2] hash mode. This hash mode which follows the message input to hash function should have to be same as the plain text input to the block cipher.

# 2.2.3. Optional Arguments

To increase the flexibility of Skein, Some Optional Inputs can be enabled. The key which should process first to achieve the security. But is optional. KEY:

### **CONFIGURATION:**

It is required in order to encode the desired output length.

### **PERSONALIZATION:**

It is used to create different functions for different users.

### PUBLIC KEY:

When hashing a message for signing it is must to public key.

#### **KEY DERIVATION IDENTIFIER:**

It is used for key derivation. To derive a key, master key should be given as input key and the requested identifier will get the key.

### NONCE:

It is particularly used in Stream cipher mode and randomized hashing [1].

### **MESSAGE:**

It is the normal message given as the Input to the hash function.

#### **OUTPUT**:

It is the one which is required and retrieved. But the main advantage of this hashing function is that whenever in need it is possible to add any function which is standardized here.

# III. Skein Hash architecture

It has different types of architecture. It will face either in Hardware or software mode. The skein

hash algorithm's structured and its different types of architecture will be discussed here.

#### 3.1. Algorithm



Fig: 4.Stick diagram



#### 3.2. Architecture

This is the architecture which is meant and defined for rounding. There are three different types as Iterative round, Loop unrolling and Pipeline architecture.

#### 3.2.1. Iterative Round architecture

One round of skein algorithm is implemented in the round circuit.Skien-512 requires the 72 number of rounds. Hence 72 times of iteration is needed to complete one hardware round for skein. The round operation not only carried by input. But also based on key, Tweak and round Number to produce a round key. The result is stored in register. If it is final output it gets XORed with input which forwards to UBI. If it is not final result then it will pass to loop till 72 rounds. Each rotation will carry over by one clock cycle exactly. If there are n number of rounds then n clock cycles are require.



Fig.6.Iterative round architecture

#### 3.2.2. Loop unrolling architecture

The main different of this architecture when compared with iterative is here the exact number of rounds have to place before registering the output. The area requirement is more due to more number of rounds placed one after another. The throughput is also low and hence lowers efficiency.



Fig:7. Loop unrolling architecture

#### **3.2.3.** Pipeline architecture

The main advantage of pipeline architecture is increasing throughput. The advantageous feature of

pipeline architecture is that each and every consecutive rounds have one register.

Skien is the type of hashing algorithm which enforces data serialization.Hashing block of n meaasge which cannot continue if n-1 value is not there.The drawback of this is that it is ideal for some rounds and hence inefficient utilization of hardware resources



Fig:8. Pipeline architecture

Messages												
Itb/clk	1											1
🔶 /tb/rst	0											
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# **IV. Experimental results**

#### MIX FUNCTION



### Key schedule



### V. Conclusion

Thus in this paper which I discussed the basic building blocks of skein hash algorithm and its architecture with detail manner. Skein has three different internal state sizes as 256, 512, 1024 bits and output may be any size. It is a good replacement of SHA family. Skein is an efficient tool to be used for large number of functions due to its optional and extendable argument system [1]. Skein is efficient for both hardware and software platforms.

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